Script generated by TTT

Title: Seidl: Virtual Machines (05.05.2014)

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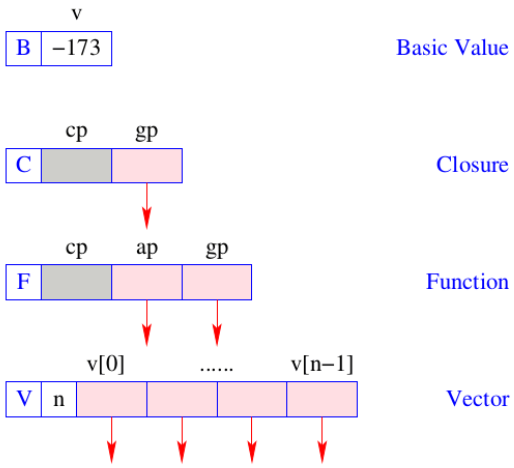
Pages: 43

... it can be thought of as an abstract data type, being capable of holding data objects of the following form:



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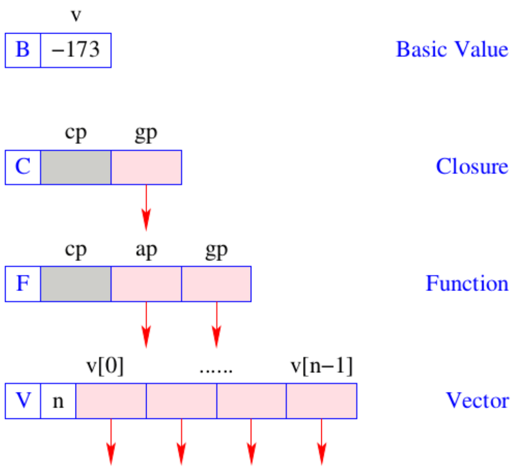

The instruction new $(tag, args)$ creates a corresponding object (B, C, F, V) in H and returns a reference to it.

We distinguish three different kinds of code for an expression e:

- $\text{code}_V\ e$ — (generates code that) computes the Value of e, stores it in the heap and returns a reference to it on top of the stack (the normal case);
- $\text{code}_B\ e$ — computes the value of e, and returns it on the top of the stack (only for Basic types);
- $\text{code}_C\ e$ — does not evaluate e, but stores a Closure of e in the heap and returns a reference to the closure on top of the stack.

We start with the code schemata for the first two kinds:

13 Simple expressions

Expressions consisting only of constants, operator applications, and conditionals are translated like expressions in imperative languages:

$$\begin{aligned}
\text{code}_B\; b\; \rho\; \text{sd} &= \text{loadc b} \\
\text{code}_B\; (\Box_1\; e)\; \rho\; \text{sd} &= \text{code}_B\; e\; \rho\; \text{sd} \\
&\; \text{op}_1 \\
\text{code}_B\; (e_1\; \Box_2\; e_2)\; \rho\; \text{sd} &= \text{code}_B\; e_1\; \rho\; \text{sd} \\
&\; \text{code}_B\; e_2\; \rho\; (\text{sd}+1) \\
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\end{aligned}$$

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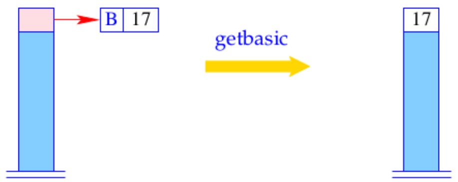
$$\begin{aligned}
\text{code}_B\; (\textbf{if}\; e_0\; \textbf{then}\; e_1\; \textbf{else}\; e_2)\; \rho\; \text{sd} \quad = &\quad \text{code}_B\; e_0\; \rho\; \text{sd} \\
&\quad \text{jumpz A} \\
&\quad \text{code}_B\; e_1\; \rho\; \text{sd} \\
&\quad \text{jump B} \\
\text{A:} &\quad \text{code}_B\; e_2\; \rho\; \text{sd} \\
\text{B:} &\quad \ldots
\end{aligned}$$

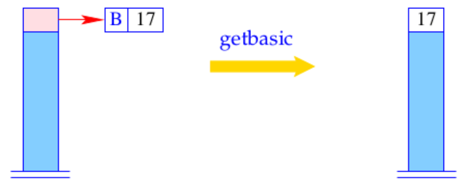
Note:

- ρ denotes the actual address environment, in which the expression is translated.
- The extra argument sd, the stack difference, *simulates* the movement of the SP when instruction execution modifies the stack. It is needed later to address variables.
- The instructions op_1 and op_2 implement the operators $\Box_1$ and $\Box_2$, in the same way as the the operators neg and add implement negation resp. addition in the CMa.
- For all other expressions, we first compute the value in the heap and then dereference the returned pointer:

$$\begin{aligned}
\text{code}_B\; e\; \rho\; \text{sd} &= \text{code}_V\; e\; \rho\; \text{sd} \\
&\; \text{getbasic}
\end{aligned}$$




```
if (H[S[SP]] != (B,_))
    Error “not basic!”;
else
    S[SP] = H[S[SP]].v;
```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For code_V and simple expressions, we define analogously:

$$
\begin{array}{lcll}
\text{code}_V\ b\ \rho\ \text{sd} & = & & \text{loadc } b\text{; mkbasic} \\
\text{code}_V\ (\Box_1\ e)\ \rho\ \text{sd} & = & & \text{code}_B\ e\ \rho\ \text{sd} \\
 & & & \text{op}_1\text{; mkbasic} \\
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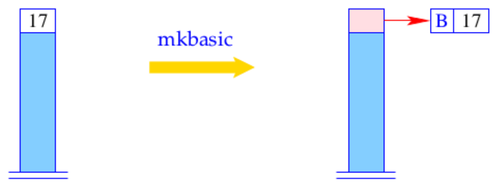
 & & & \text{code}_B\ e_2\ \rho\ (\text{sd}+1) \\
 & & & \text{op}_2\text{; mkbasic} \\
\text{code}_V\ (\textbf{if}\ e_0\ \textbf{then}\ e_1\ \textbf{else}\ e_2)\ \rho\ \text{sd} & = & & \text{code}_B\ e_0\ \rho\ \text{sd} \\
 & & & \text{jumpz A} \\
 & & & \text{code}_V\ e_1\ \rho\ \text{sd} \\
 & & & \text{jump B} \\
 & & \text{A:} & \text{code}_V\ e_2\ \rho\ \text{sd} \\
 & & \text{B:} & \ldots
\end{array}
$$




```
S[SP] = new (B,S[SP]);
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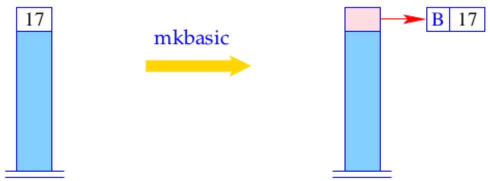
 & & A: & \mathrm{code}_V\; e_2\; \rho\; \mathrm{sd} \\
 & & B: & \ldots
\end{array}
$$



S[SP] = new (B,S[SP]);

14 Accessing Variables

We must distinguish between local and global variables.

Example: Regard the function f :

$$
\begin{array}{lllll}
\mathbf{let} & c = 5 & & & \\
\mathbf{in\ let} & f = \mathbf{fun}\; a & \rightarrow & \mathbf{let}\; b = a * a & \\
 & & & \mathbf{in}\; b + c & \\
\mathbf{in} & f\; c & & &
\end{array}
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The function f uses the global variable c and the local variables a (as formal parameter) and b (introduced by the inner **let**).

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Accessing Global Variables

- The bindings of global variables of an expression or a function are kept in a vector in the heap (Global Vector).
- They are addressed consecutively starting with 0.
- When an F-object or a C-object are constructed, the Global Vector for the function or the expression is determined and a reference to it is stored in the gp-component of the object.
- During the evaluation of an expression, the (new) register GP (Global Pointer) points to the actual Global Vector.
- In constrast, local variables should be administered on the stack ...

$\Longrightarrow$ General form of the address environment:

$$\rho : Vars \to \{L, G\} \times \mathbb{Z}$$

Accessing Local Variables

Local variables are administered on the stack, in stack frames.

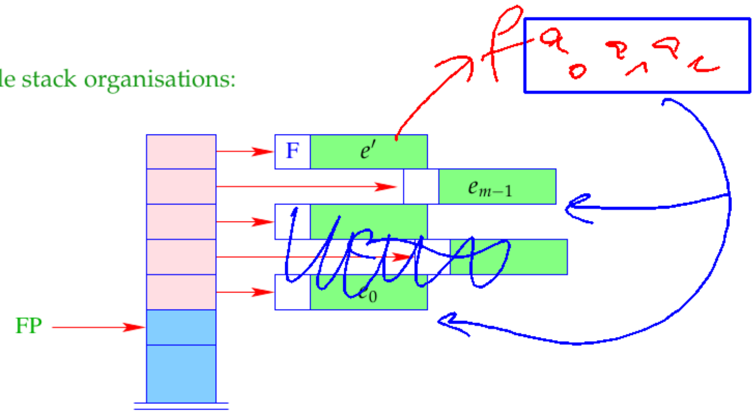
Let $e \equiv e' \; e_0 \; \ldots \; e_{m-1}$ be the application of a function e' to arguments $e_0, \ldots, e_{m-1}$.

Warning:

The arity of e' does not need to be m :-)

- f may therefore receive less than n arguments (under supply);
- f may also receive more than n arguments, if t is a functional type (over supply).

Possible stack organisations:

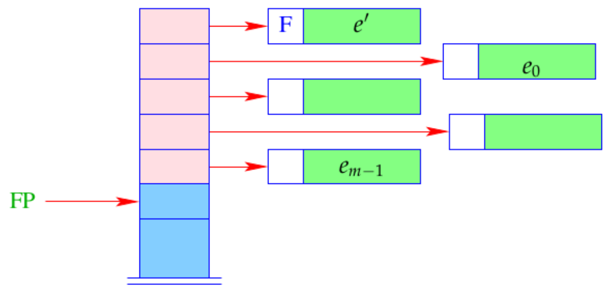


\+ Addressing of the arguments can be done relative to FP

– The local variables of e' cannot be addressed relative to FP.

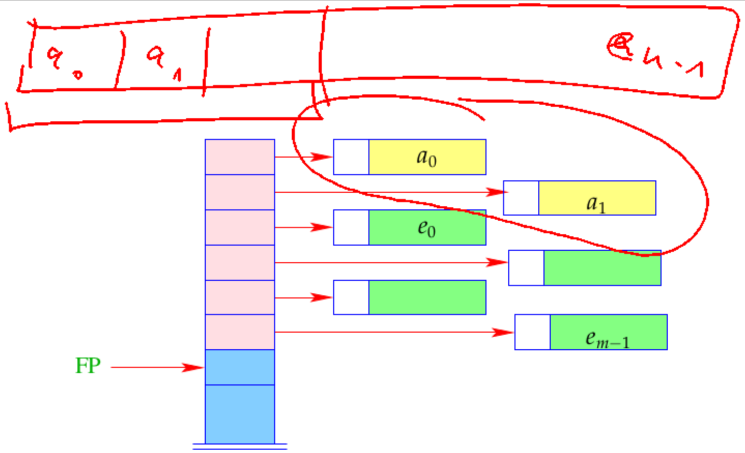
– If e' is an n-ary function with $n < m$, i.e., we have an over-supplied function application, the remaining $m - n$ arguments will have to be shifted.

Alternative:



\+ The further arguments $a_0, \ldots, a_{k-1}$ and the local variables can be allocated above the arguments.

- Addressing of arguments and local variables relative to FP is no more possible. (Remember: m is unknown when the function definition is translated.)



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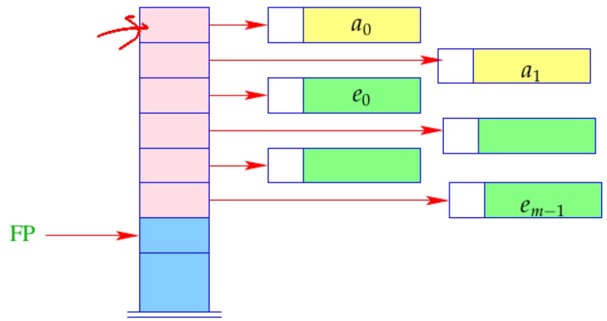

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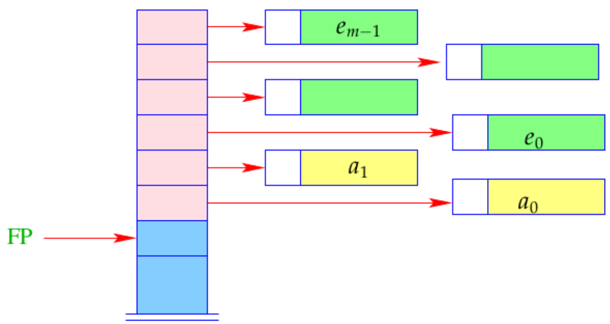
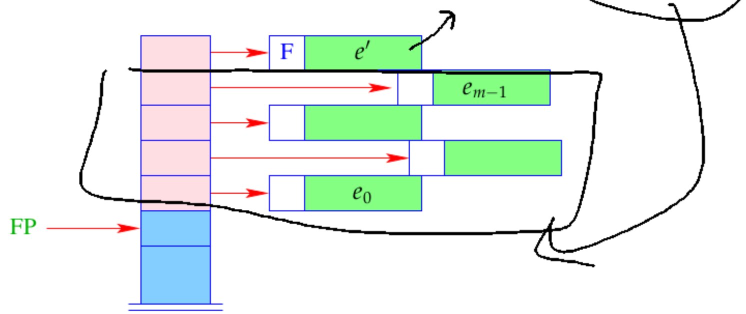
- If e' evaluates to a function, which has already been partially applied to the parameters $a_0, \dots, a_{k-1}$, these have to be sneaked in underneath e_0:



Possible stack organisations: let $g = f(a_0, a_1, a_2)$



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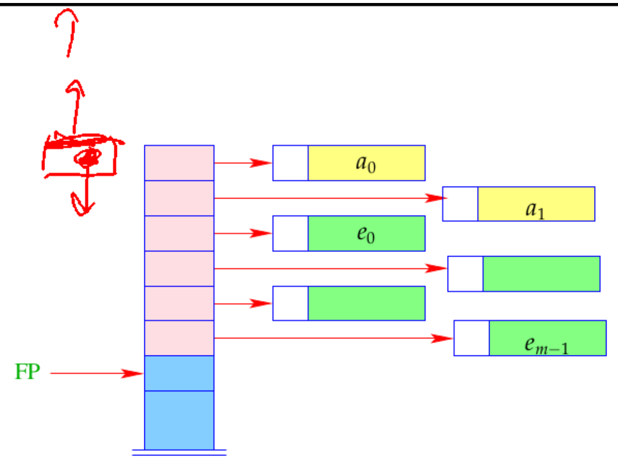

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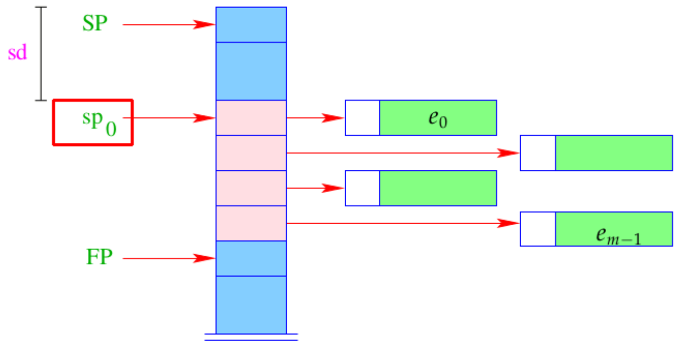
Way out:

- We address both, arguments and local variables, relative to the stack pointer SP !!!
- However, the stack pointer changes during program execution...



- The differerence between the current value of SP and its value sp_0 at the entry of the function body is called the stack distance, sd.
- Fortunately, this stack distance can be determined at compile time for each program point, by simulating the movement of the SP.
- The formal parameters $x_0, x_1, x_2, \ldots$ successively receive the non-positive relative addresses $0, -1, -2, \ldots$, i.e., $\rho\, x_i = (L, -i)$.
- The absolute address of the i-th formal parameter consequently is

$$\text{sp}_0 - i = (\text{SP} - \text{sd}) - i$$

- The local **let**-variables $y_1, y_2, y_3, \ldots$ will be successively pushed onto the stack:

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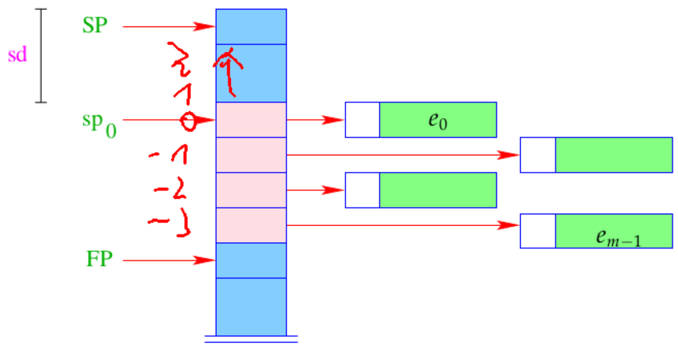

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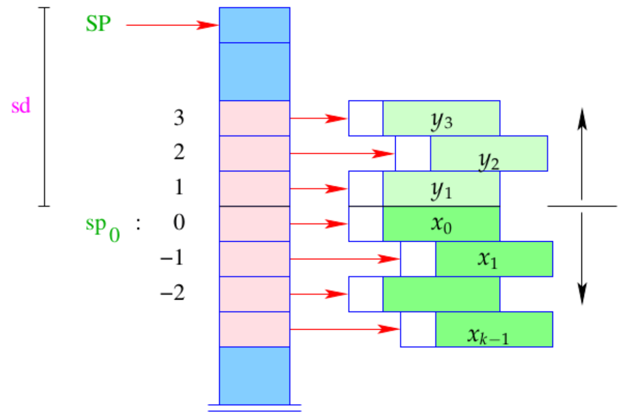


- The y_i have positive relative addresses $1, 2, 3, \ldots$, that is: $\rho\; y_i = (L, i)$.
- The absolute address of y_i is then $\mathrm{sp}_0 + i = (\mathrm{SP} - \mathrm{sd}) + i$

With CBN, we generate for the access to a variable:

$$\begin{array}{rcl} \mathrm{code}_V\; x\; \rho\; \mathrm{sd} & = & \mathrm{getvar}\; x\; \rho\; \mathrm{sd} \\ & & \mathrm{eval} \end{array}$$

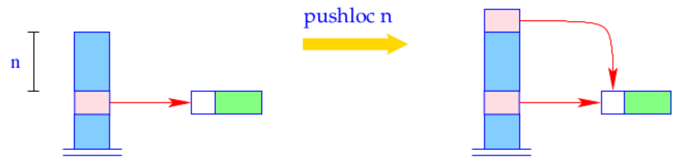
The instruction eval checks, whether the value has already been computed or whether its evaluation has to yet to be done ($\Longrightarrow$ will be treated later :-)

With CBV, we can just delete eval from the above code schema.

The (compile-time) macro getvar is defined by:

$$\begin{array}{rcl} \mathrm{getvar}\; x\; \rho\; \mathrm{sd} & = & \mathbf{let}\; (t, i) = \rho\; x\; \mathbf{in} \\ & & \mathbf{match}\; t\; \mathbf{with} \\ & & \quad L \to \mathrm{pushloc}\; (\mathrm{sd} - i) \\ & & \quad \mid G \to \mathrm{pushglob}\; \mathrm{i} \\ & & \mathbf{end} \end{array}$$

The access to local variables:



S[SP+1] =S[SP - n]; SP++;

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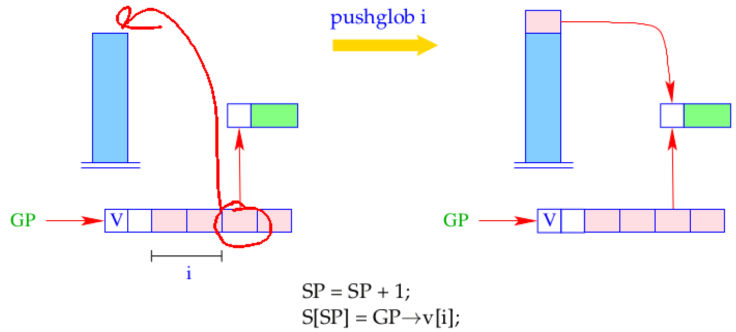
$SP - (Sd - i) =$
$(SP - SD) + i$

The access to global variables is much simpler:



Correctness argument:

Let sp and sd be the values of the stack pointer resp. stack distance before the execution of the instruction. The value of the local variable with address i is loaded from $S[a]$ with

$$a = \text{sp} - (\text{sd} - i) = (\text{sp} - \text{sd}) + i = \text{sp}_0 + i$$

... exactly as it should be :-)

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Example:

Regard $e \equiv (b + c)$ for $\rho = \{b \mapsto (L, 1), c \mapsto (G, 0)\}$ and sd $= 1$.

With CBN, we obtain:

$$
\begin{array}{lllll}
\text{code}_V\; e\; \rho\; 1 & = & \boxed{\text{getvar}\; b\; \rho\; 1} & = & 1 \quad \boxed{\text{pushloc}\; 0} \\
 & & \text{eval} & & 2 \quad \cancel{\text{eval}} \\
 & & \text{getbasic} & & 2 \quad \text{getbasic} \\
 & & \boxed{\text{getvar}\; c\; \rho\; 2} & & 2 \quad \boxed{\text{pushglob}\; 0} \\
 & & \text{eval} & & 3 \quad \cancel{\text{eval}} \\
 & & \text{getbasic} & & 3 \quad \text{getbasic} \\
 & & \text{add} & & 3 \quad \text{add} \\
 & & \text{mkbasic} & & 2 \quad \text{mkbasic}
\end{array}
$$

15 let-Expressions

As a warm-up let us first consider the treatment of local variables :-)

Let $e \equiv \textbf{let}\; y_1 = e_1\; \textbf{in} \ldots \textbf{let}\; e_n\; \textbf{in}\; e_0$ be a nested **let**-expression.

The translation of e must deliver an instruction sequence that

- allocates local variables $y_1, \ldots, y_n$;
- in the case of
  CBV: evaluates $e_1, \ldots, e_n$ and binds the y_i to their values;
  CBN: constructs closures for the $e_1, \ldots, e_n$ and binds the y_i to them;
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Here, we consider the non-recursive case only, i.e. where y_j only depends on $y_1, \ldots, y_{j-1}$. We obtain for CBN:

15 let-Expressions

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\text{code}_V\ e\ \rho\ \text{sd} \quad = \quad & \text{code}_C\ e_1\ \rho\ \text{sd} \\
& \text{code}_C\ e_2\ \rho_1\ (\text{sd}+1) \\
& \ldots \\
& \text{code}_C\ e_n\ \rho_{n-1}\ (\text{sd}+n-1) \\
& \text{code}_V\ e_0\ \rho_n\ (\text{sd}+n) \\
& \text{slide n} \qquad\qquad\qquad // \text{ deallocates local variables}
\end{aligned}$$

where $\quad \rho_j = \rho \oplus \{y_i \mapsto (L, \text{sd}+i) \mid i = 1, \ldots, j\}$.

In the case of CBV, we use code_V for the expressions $e_1, \ldots, e_n$.

Warning!

All the e_i must be associated with the same binding for the global variables!

$a \mapsto (L, 1)$
$b \mapsto (L, 2)$

Example:

Consider the expression

$$e \equiv \textbf{let } a = 19 \textbf{ in let } b = a * a \textbf{ in } a + b$$

for $\rho = \emptyset$ and sd $= 0$. We obtain (for CBV):

0	loadc 19	3	getbasic	3	pushloc 1
1	mkbasic	3	mul	4	getbasic
1	pushloc 0	2	mkbasic	4	add
2	getbasic	2	pushloc 1	3	mkbasic
2	pushloc 1	3	getbasic	3	slide 2

$$\begin{aligned}
\text{code}_V\ e\ \rho\ \text{sd} \quad = \quad & \text{code}_C\ e_1\ \rho\ \text{sd} \\
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& \ldots \\
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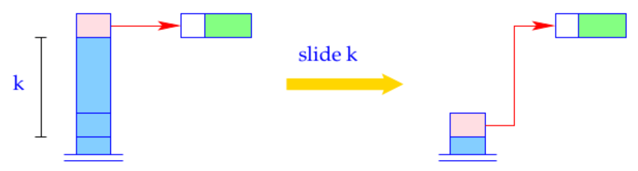
The instruction slide k deallocates again the space for the locals:



S[SP-k] = S[SP];
SP = SP - k;