Script **generated by TTT**

Title: Seidl: Programmoptimierung (14.11.2012)

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1.5 Interval Analysis

Observation:

- Programmers often use global constants for switching debugging code on/off.

  $\Longrightarrow$

  Constant propagation is useful :-)

- In general, precise values of variables will be unknown — perhaps, however, a tight interval !!!

Example:

```
for (i = 0; i < 42; i++)
    if (0 ≤ i ∧ i < 42){
        A1 = A + i;
        M[A1] = i;
    }
// A start address of an array
// if the array-bound check
```

Obviously, the inner check is superfluous :-)

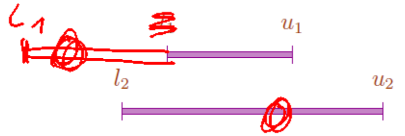
Idea 1:

Determine for every variable x an (as tight as possible :-) interval of possible values:

$$\mathbb{I} = \{[l, u] \mid l \in \mathbb{Z} \cup \{-\infty\}, u \in \mathbb{Z} \cup \{+\infty\}, l \leq u\}$$

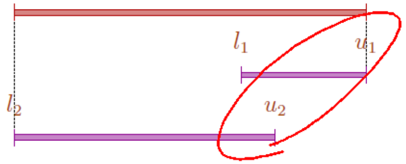
Partial Ordering:

$$[l_1, u_1] \sqsubseteq [l_2, u_2] \quad \text{iff} \quad l_2 \leq l_1 \wedge u_1 \leq u_2$$



Thus:

$$[l_1, u_1] \sqcup [l_2, u_2] \quad = \quad [l_1 \sqcap l_2, u_1 \sqcup u_2]$$



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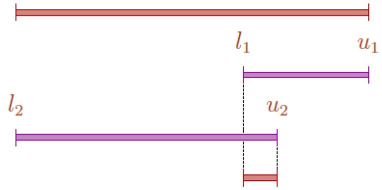
$$[l_1, u_1] \sqcup [l_2, u_2] \quad = \quad [l_1 \sqcap l_2, u_1 \sqcup u_2]$$

$$[l_1, u_1] \sqcap [l_2, u_2] \quad = \quad [l_1 \sqcup l_2, u_1 \sqcap u_2] \qquad \text{whenever } (l_1 \sqcup l_2) \leq (u_1 \sqcap u_2)$$



Caveat:

→ $\mathbb{I}$ is not a complete lattice :-)

→ $\mathbb{I}$ has infinite ascending chains, e.g.,

$$[0,0] \sqsubset [0,1] \sqsubset [-1,1] \sqsubset [-1,2] \sqsubset \ldots$$

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Description Relation:

$$z \; \Delta \; [l, u] \qquad \text{iff} \qquad l \leq z \leq u$$

Concretization:

$$\gamma \, [l, u] = \{z \in \mathbb{Z} \mid l \leq z \leq u\}$$

Example:

$$\gamma\,[0,7] = \{0,\ldots,7\}$$
$$\gamma\,[0,\infty] = \{0,1,2,\ldots,\}$$

Computing with intervals: Interval Arithmetic :-)

Addition:

$$[l_1,u_1] +^\sharp [l_2,u_2] = [l_1+l_2, u_1+u_2] \quad \text{where}$$
$$-\infty + \_ = -\infty$$
$$+\infty + \_ = +\infty$$
$$// \; -\infty + \infty \quad \text{cannot occur} \quad \text{:-)}$$

Negation:

$$-^\sharp\,[l,u] = [-u,-l]$$

Multiplication:

$$[l_1,u_1] *^\sharp [l_2,u_2] = [a,b] \quad \text{where}$$
$$a = l_1l_2 \sqcap l_1u_2 \sqcap u_1l_2 \sqcap u_1u_2$$
$$b = l_1l_2 \sqcup l_1u_2 \sqcup u_1l_2 \sqcup u_1u_2$$

Example:

$$[0,2] *^\sharp [3,4] = [0,8]$$
$$[-1,2] *^\sharp [3,4] = [-4,8]$$
$$[-1,2] *^\sharp [-3,4] = [-6,8]$$
$$[-1,2] *^\sharp [-4,-3] = [-8,4]$$

Division:

$$[l_1,u_1] /^\sharp [l_2,u_2] = [a,b]$$

- If 0 is not contained in the interval of the denominator, then:

$$a = l_1/l_2 \sqcap l_1/u_2 \sqcap u_1/l_2 \sqcap u_1/u_2$$
$$b = l_1/l_2 \sqcup l_1/u_2 \sqcup u_1/l_2 \sqcup u_1/u_2$$

- If: $l_2 \leq 0 \leq u_2$ we define:

$$[a,b] = [-\infty,+\infty]$$

Equality:

$$[l_1,u_1] ==^\sharp [l_2,u_2] = \begin{cases} [1,1] & \text{if} \quad l_1 = u_1 = l_2 = u_2 \\ [0,0] & \text{if} \quad u_1 < l_2 \vee u_2 < l_1 \\ [0,1] & \text{otherwise} \end{cases}$$

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Example:

$$\begin{array}{rcl} [42,42] ==^\sharp [42,42] & = & [1,1] \\ [0,7] ==^\sharp [0,7] & = & [0,1] \\ [1,2] ==^\sharp [3,4] & = & [0,0] \end{array}$$

Less:

$$[l_1,u_1] <^\sharp [l_2,u_2] \;=\; \begin{cases} [1,1] & \text{if} \quad u_1 < l_2 \\ [0,0] & \text{if} \quad u_2 \leq l_1 \\ [0,1] & \text{otherwise} \end{cases}$$

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Example:

$$\begin{array}{rcl} [1,2] <^\sharp [9,42] & = & [1,1] \\ [0,7] <^\sharp [0,7] & = & [0,1] \\ [3,4] <^\sharp [1,2] & = & [0,0] \end{array}$$

By means of $\mathbb{I}$ we construct the complete lattice:

$$\mathbb{D}_{\mathbb{I}} = (\mathit{Vars} \to \mathbb{I})_\perp$$

Description Relation:

$$\rho \;\Delta\; D \qquad \text{iff} \qquad D \neq \bot \quad \wedge \quad \forall x \in \mathit{Vars} : (\rho\, x) \;\Delta\; (D\, x)$$

The abstract evaluation of expressions is defined analogously to constant propagation. We have:

$$([\![e]\!]\, \rho) \;\Delta\; ([\![e]\!]^\sharp\, D) \qquad \text{whenever} \qquad \rho \;\Delta\; D$$

The Effects of Edges:

$$
\begin{array}{lcl}
[\![;]\!]^\sharp\, D & = & D \\
[\![x = e;]\!]^\sharp\, D & = & D \oplus \{x \mapsto [\![e]\!]^\sharp\, D\} \\
[\![x = M[e];]\!]^\sharp\, D & = & D \oplus \{x \mapsto \top\} \\
[\![M[e_1] = e_2;]\!]^\sharp\, D & = & D \\
[\![\text{Pos}\,(e)]\!]^\sharp\, D & = & \begin{cases} \bot & \text{if} \quad [0,0] = [\![e]\!]^\sharp\, D \\ D & \text{otherwise} \end{cases} \\
[\![\text{Neg}\,(e)]\!]^\sharp\, D & = & \begin{cases} D & \text{if} \quad [0,0] \sqsubseteq [\![e]\!]^\sharp\, D \\ \bot & \text{otherwise} \end{cases}
\end{array}
$$

... given that $\quad D \neq \bot \quad$:-)

Better Exploitation of Conditions:

$$
[\![\text{Pos}\,(e)]\!]^\sharp\, D \;=\; \begin{cases} \bot & \text{if} \quad [0,0] = [\![e]\!]^\sharp\, D \\ D_1 & \text{otherwise} \end{cases}
$$

where :

$$
D_1 \;=\; \begin{cases} D \oplus \{x \mapsto (D\,x) \sqcap ([\![e_1]\!]^\sharp\, D)\} & \text{if } e \equiv x == e_1 \\ D \oplus \{x \mapsto (D\,x) \sqcap [-\infty, u]\} & \text{if } e \equiv x \leq e_1, [\![e_1]\!]^\sharp\, D = [\_, u] \\ D \oplus \{x \mapsto (D\,x) \sqcap [l, \infty]\} & \text{if } e \equiv x \geq e_1, [\![e_1]\!]^\sharp\, D = [l, \_] \end{cases}
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$$

x == e

x < y

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332

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333

Better Exploitation of Conditions (cont.):

$$[\![\mathrm{Neg}\,(e)]\!]^\sharp\, D = \begin{cases} \bot & \text{if } [0,0] \not\sqsubseteq [\![e]\!]^\sharp\, D \\ D_1 & \text{otherwise} \end{cases}$$

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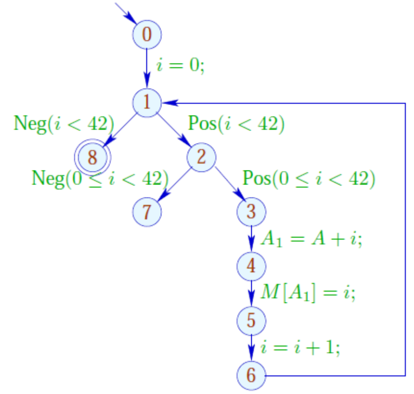

	i	
	l	u
0	$-\infty$	$+\infty$
1	0	42
2	0	41
3	0	41
4	0	41
5	0	41
6	1	42
7	$\bot$	
8	42	42

Problem:

→ The solution can be computed with RR-iteration —
after about 42 rounds :-(

→ On some programs, iteration may never terminate :-((

Idea 1: Widening

- Accelerate the iteration — at the prize of imprecision :-)
- Allow only a bounded number of modifications of values !!!

... in the Example:

- dis-allow updates of interval bounds in $\mathbb{Z}$ ...

$\Longrightarrow$ a maximal chain:

$$[3, 17] \sqsubset [3, +\infty] \sqsubset [-\infty, +\infty]$$